

The theory of Steiner triple systems

Silvia Barbina

joint work with Enrique Casanovas

LC2018, Udine

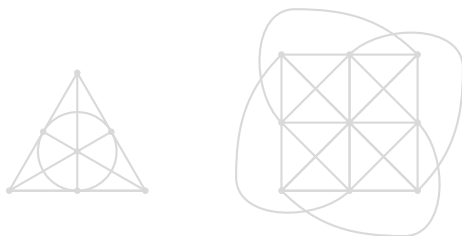
Steiner triple systems

Definition

A finite **Steiner triple system** (STS) of order n is a pair (V, \mathcal{B}) where:

- V is a set of n elements;
- \mathcal{B} is a collection of 3-element subsets of V (the **blocks**) such that any two $x, y \in V$ are contained in exactly one block.

A set V with a collection of 3-element subsets is a **partial STS** if any two elements of V belong to at most one block.



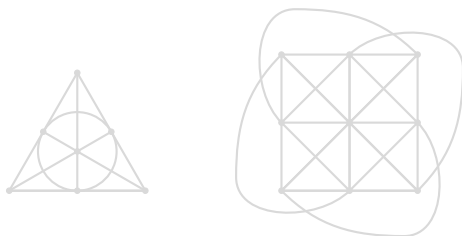
Steiner triple systems

Definition

A finite **Steiner triple system** (STS) of order n is a pair (V, \mathcal{B}) where:

- V is a set of n elements;
- \mathcal{B} is a collection of 3-element subsets of V (the **blocks**) such that any two $x, y \in V$ are contained in exactly one block.

A set V with a collection of 3-element subsets is a **partial STS** if any two elements of V belong to at most one block.



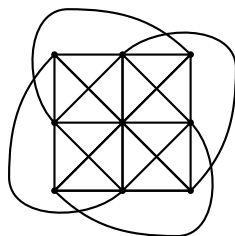
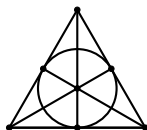
Steiner triple systems

Definition

A finite **Steiner triple system** (STS) of order n is a pair (V, \mathcal{B}) where:

- V is a set of n elements;
- \mathcal{B} is a collection of 3-element subsets of V (the **blocks**) such that any two $x, y \in V$ are contained in exactly one block.

A set V with a collection of 3-element subsets is a **partial STS** if any two elements of V belong to at most one block.



Kirkman's schoolgirl problem

Fifteen girls in a school take a walk in rows of three for seven days in succession. Is there an arrangement such that no two girls walk together in a row more than once?

(Thomas Penyngton Kirkman, 1850)

STSs appear in

- combinatorial design theory (they are balanced incomplete block designs)
- design of experiments
- coding theory.

More general Steiner systems are connected to the Mathieu groups.

Kirkman's schoolgirl problem

Fifteen girls in a school take a walk in rows of three for seven days in succession. Is there an arrangement such that no two girls walk together in a row more than once?

(Thomas Penyngton Kirkman, 1850)

STSs appear in

- combinatorial design theory (they are balanced incomplete block designs)
- design of experiments
- coding theory.

More general Steiner systems are connected to the Mathieu groups.

Kirkman's schoolgirl problem

Fifteen girls in a school take a walk in rows of three for seven days in succession. Is there an arrangement such that no two girls walk together in a row more than once?

(Thomas Penyngton Kirkman, 1850)

STSs appear in

- combinatorial design theory (they are balanced incomplete block designs)
- design of experiments
- coding theory.

More general Steiner systems are connected to the Mathieu groups.

Kirkman's schoolgirl problem

Fifteen girls in a school take a walk in rows of three for seven days in succession. Is there an arrangement such that no two girls walk together in a row more than once?

(Thomas Penyngton Kirkman, 1850)

STSs appear in

- combinatorial design theory (they are balanced incomplete block designs)
- design of experiments
- coding theory.

More general Steiner systems are connected to the Mathieu groups.

Kirkman's schoolgirl problem

Fifteen girls in a school take a walk in rows of three for seven days in succession. Is there an arrangement such that no two girls walk together in a row more than once?

(Thomas Penyngton Kirkman, 1850)

STSs appear in

- combinatorial design theory (they are balanced incomplete block designs)
- design of experiments
- coding theory.

More general Steiner systems are connected to the Mathieu groups.

Steiner triple systems

- When n is finite, an STS of order n exists if and only if $n \equiv 1$ or $3 \pmod{6}$.
- If we allow $|V| \geq \omega$, the pair (V, \mathcal{B}) is an **infinite STS**.

We can describe blocks via

- a ternary relation R where $R(x, y, z)$ if and only if $\{x, y, z\}$ is a block, or
- a binary operation \cdot defined by

$$x \cdot y = z \text{ iff } \{x, y, z\} \text{ is a block.}$$

When blocks are described by a relation, a substructure of an STS is a *partial* STS.

In a functional language, substructures are STSs.

Steiner triple systems

- When n is finite, an STS of order n exists if and only if $n \equiv 1$ or $3 \pmod{6}$.
- If we allow $|V| \geq \omega$, the pair (V, \mathcal{B}) is an **infinite STS**.

We can describe blocks via

- a ternary relation R where $R(x, y, z)$ if and only if $\{x, y, z\}$ is a block, or
- a binary operation \cdot defined by

$$x \cdot y = z \text{ iff } \{x, y, z\} \text{ is a block.}$$

When blocks are described by a relation, a substructure of an STS is a *partial* STS.

In a functional language, substructures are STSs.

Steiner triple systems

- When n is finite, an STS of order n exists if and only if $n \equiv 1$ or $3 \pmod{6}$.
- If we allow $|V| \geq \omega$, the pair (V, \mathcal{B}) is an **infinite STS**.

We can describe blocks via

- a ternary relation R where $R(x, y, z)$ if and only if $\{x, y, z\}$ is a block, or
- a binary operation \cdot defined by

$$x \cdot y = z \text{ iff } \{x, y, z\} \text{ is a block.}$$

When blocks are described by a relation, a substructure of an STS is a *partial* STS.

In a functional language, substructures are STSs.

Steiner triple systems

- When n is finite, an STS of order n exists if and only if $n \equiv 1$ or $3 \pmod{6}$.
- If we allow $|V| \geq \omega$, the pair (V, \mathcal{B}) is an **infinite STS**.

We can describe blocks via

- a ternary relation R where $R(x, y, z)$ if and only if $\{x, y, z\}$ is a block, or
- a binary operation \cdot defined by

$$x \cdot y = z \text{ iff } \{x, y, z\} \text{ is a block.}$$

When blocks are described by a relation, a substructure of an STS is a *partial* STS.

In a functional language, substructures are STSs.

Steiner triple systems

- When n is finite, an STS of order n exists if and only if $n \equiv 1$ or $3 \pmod{6}$.
- If we allow $|V| \geq \omega$, the pair (V, \mathcal{B}) is an **infinite STS**.

We can describe blocks via

- a ternary relation R where $R(x, y, z)$ if and only if $\{x, y, z\}$ is a block, or
- a binary operation \cdot defined by

$$x \cdot y = z \text{ iff } \{x, y, z\} \text{ is a block.}$$

When blocks are described by a relation, a substructure of an STS is a *partial* STS.

In a functional language, substructures are STSs.

Steiner triple systems

- When n is finite, an STS of order n exists if and only if $n \equiv 1$ or $3 \pmod{6}$.
- If we allow $|V| \geq \omega$, the pair (V, \mathcal{B}) is an **infinite STS**.

We can describe blocks via

- a ternary relation R where $R(x, y, z)$ if and only if $\{x, y, z\}$ is a block, or
- a binary operation \cdot defined by

$$x \cdot y = z \text{ iff } \{x, y, z\} \text{ is a block.}$$

When blocks are described by a relation, a substructure of an STS is a *partial* STS.

In a functional language, substructures are STSs.

Steiner triple systems

- When n is finite, an STS of order n exists if and only if $n \equiv 1$ or $3 \pmod{6}$.
- If we allow $|V| \geq \omega$, the pair (V, \mathcal{B}) is an **infinite STS**.

We can describe blocks via

- a ternary relation R where $R(x, y, z)$ if and only if $\{x, y, z\}$ is a block, or
- a binary operation \cdot defined by

$$x \cdot y = z \text{ iff } \{x, y, z\} \text{ is a block.}$$

When blocks are described by a relation, a substructure of an STS is a *partial* STS.

In a functional language, substructures are STSs.

We choose a functional language, so that an STS is a structure (A, \cdot) where \cdot is a binary operation on A such that

- 1 $x \cdot y = y \cdot x$
- 2 $x \cdot x = x$
- 3 $x \cdot (x \cdot y) = y.$

Definition

T_{STS} is the theory that contains axioms 1–3 above.

T_{STS} is a universal theory.

We choose a functional language, so that an STS is a structure (A, \cdot) where \cdot is a binary operation on A such that

- 1 $x \cdot y = y \cdot x$
- 2 $x \cdot x = x$
- 3 $x \cdot (x \cdot y) = y.$

Definition

T_{STS} is the theory that contains axioms 1–3 above.

T_{STS} is a universal theory.

We choose a functional language, so that an STS is a structure (A, \cdot) where \cdot is a binary operation on A such that

- 1 $x \cdot y = y \cdot x$
- 2 $x \cdot x = x$
- 3 $x \cdot (x \cdot y) = y.$

Definition

T_{STS} is the theory that contains axioms 1–3 above.

T_{STS} is a universal theory.

Fact

- 1 *Every finite partial STS can be embedded in a finite STS.*
- 2 *Every infinite partial STS can be embedded in an STS of the same cardinality.*

Fact

- 1 *Every finite partial STS can be embedded in a finite STS.*
- 2 *Every infinite partial STS can be embedded in an STS of the same cardinality.*

The Fraïssé limit of the finite Steiner triple systems

The class \mathcal{C} of all finite Steiner triple systems has

- the Joint Embedding and the Amalgamation Properties
- the Hereditary Property
- countably many isomorphism types.

Therefore \mathcal{C} has a Fraïssé limit: the unique (up to isomorphism) countable Steiner triple system M_F which is *ultrahomogeneous* and *universal* (for finite Steiner triple systems).

Questions

*What can we say about $\text{Th}(M_F)$? Can we describe its models?
Does it have q.e.?*

The Fraïssé limit of the finite Steiner triple systems

The class \mathcal{C} of all finite Steiner triple systems has

- the Joint Embedding and the Amalgamation Properties
- the Hereditary Property
- countably many isomorphism types.

Therefore \mathcal{C} has a Fraïssé limit: the unique (up to isomorphism) countable Steiner triple system M_F which is *ultrahomogeneous* and *universal* (for finite Steiner triple systems).

Questions

*What can we say about $\text{Th}(M_F)$? Can we describe its models?
Does it have q.e.?*

The Fraïssé limit of the finite Steiner triple systems

The class \mathcal{C} of all finite Steiner triple systems has

- the Joint Embedding and the Amalgamation Properties
- the Hereditary Property
- countably many isomorphism types.

Therefore \mathcal{C} has a Fraïssé limit: the unique (up to isomorphism) countable Steiner triple system M_F which is *ultrahomogeneous* and *universal* (for finite Steiner triple systems).

Questions

*What can we say about $\text{Th}(M_F)$? Can we describe its models?
Does it have q.e.?*

The Fraïssé limit of the finite Steiner triple systems

The class \mathcal{C} of all finite Steiner triple systems has

- the Joint Embedding and the Amalgamation Properties
- the Hereditary Property
- countably many isomorphism types.

Therefore \mathcal{C} has a Fraïssé limit: the unique (up to isomorphism) countable Steiner triple system M_F which is *ultrahomogeneous* and *universal* (for finite Steiner triple systems).

Questions

*What can we say about $\text{Th}(M_F)$? Can we describe its models?
Does it have q.e.?*

The Fraïssé limit of the finite Steiner triple systems

The class \mathcal{C} of all finite Steiner triple systems has

- the Joint Embedding and the Amalgamation Properties
- the Hereditary Property
- countably many isomorphism types.

Therefore \mathcal{C} has a Fraïssé limit: the unique (up to isomorphism) countable Steiner triple system M_F which is *ultrahomogeneous* and *universal* (for finite Steiner triple systems).

Questions

What can we say about $\text{Th}(M_F)$? Can we describe its models? Does it have q.e.?

Definition

Let B be a finite partial STS. Then

- δ_B is a formula that describes the diagram of B
- $A \subseteq B$ is **relatively closed** in B if for every $a, b \in A$ and $c \in B$, if $a \cdot b = c$ then $c \in A$.

Definition

If B is a finite partial STS and $A \subseteq B$ a relatively closed subset, then

$$\phi_{(A,B)} = \forall \bar{x} (\delta_A(\bar{x}) \rightarrow \exists \bar{y} \delta_B(\bar{x}, \bar{y})).$$

Let $\Delta = \{\phi_{(A,B)} : B \text{ is a finite partial STS and } A \subseteq B \text{ is a relatively closed subset}\}$.

Definition

Let B be a finite partial STS. Then

- δ_B is a formula that describes the diagram of B
- $A \subseteq B$ is **relatively closed** in B if for every $a, b \in A$ and $c \in B$, if $a \cdot b = c$ then $c \in A$.

Definition

If B is a finite partial STS and $A \subseteq B$ a relatively closed subset, then

$$\phi_{(A,B)} = \forall \bar{x} (\delta_A(\bar{x}) \rightarrow \exists \bar{y} \delta_B(\bar{x}, \bar{y})).$$

Let $\Delta = \{\phi_{(A,B)} : B \text{ is a finite partial STS and } A \subseteq B \text{ is a relatively closed subset}\}.$

Definition

Let B be a finite partial STS. Then

- δ_B is a formula that describes the diagram of B
- $A \subseteq B$ is **relatively closed** in B if for every $a, b \in A$ and $c \in B$, if $a \cdot b = c$ then $c \in A$.

Definition

If B is a finite partial STS and $A \subseteq B$ a relatively closed subset, then

$$\phi_{(A,B)} = \forall \bar{x} (\delta_A(\bar{x}) \rightarrow \exists \bar{y} \delta_B(\bar{x}, \bar{y})).$$

Let $\Delta = \{\phi_{(A,B)} : B \text{ is a finite partial STS and } A \subseteq B \text{ is a relatively closed subset}\}.$

Definition

Let B be a finite partial STS. Then

- δ_B is a formula that describes the diagram of B
- $A \subseteq B$ is **relatively closed** in B if for every $a, b \in A$ and $c \in B$, if $a \cdot b = c$ then $c \in A$.

Definition

If B is a finite partial STS and $A \subseteq B$ a relatively closed subset, then

$$\phi_{(A,B)} = \forall \bar{x} (\delta_A(\bar{x}) \rightarrow \exists \bar{y} \delta_B(\bar{x}, \bar{y})).$$

Let $\Delta = \{\phi_{(A,B)} : B \text{ is a finite partial STS and } A \subseteq B \text{ is a relatively closed subset}\}.$

Let $T_{\text{STS}}^* = \Delta \cup T_{\text{STS}}$.

Fact

$M_F \models T_{\text{STS}}^*$.

There is more.

Theorem

*The theory T_{STS}^**

- *axiomatises the existentially closed Steiner triple systems*
- *it is model complete*
- *is the model companion of T_{STS}*
- *is complete*
- *has quantifier elimination.*

M_F is a prime model of T_{STS}^ .*

Let $T_{\text{STS}}^* = \Delta \cup T_{\text{STS}}$.

Fact

$M_F \models T_{\text{STS}}^*$.

There is more.

Theorem

*The theory T_{STS}^**

- *axiomatises the existentially closed Steiner triple systems*
- *it is model complete*
- *is the model companion of T_{STS}*
- *is complete*
- *has quantifier elimination.*

M_F is a prime model of T_{STS}^ .*

Let $T_{\text{STS}}^* = \Delta \cup T_{\text{STS}}$.

Fact

$M_F \models T_{\text{STS}}^*$.

There is more.

Theorem

*The theory T_{STS}^**

- *axiomatises the existentially closed Steiner triple systems*
- *it is model complete*
- *is the model companion of T_{STS}*
- *is complete*
- *has quantifier elimination.*

M_F is a prime model of T_{STS}^ .*

Let $T_{\text{STS}}^* = \Delta \cup T_{\text{STS}}$.

Fact

$M_F \models T_{\text{STS}}^*$.

There is more.

Theorem

*The theory T_{STS}^**

- *axiomatises the existentially closed Steiner triple systems*
- *it is model complete*
- *is the model companion of T_{STS}*
- *is complete*
- *has quantifier elimination.*

M_F is a prime model of T_{STS}^ .*

Let $T_{\text{STS}}^* = \Delta \cup T_{\text{STS}}$.

Fact

$M_F \models T_{\text{STS}}^*$.

There is more.

Theorem

*The theory T_{STS}^**

- *axiomatises the existentially closed Steiner triple systems*
- *it is model complete*
- *is the model companion of T_{STS}*
- *is complete*
- *has quantifier elimination.*

M_F is a prime model of T_{STS}^ .*

Let $T_{\text{STS}}^* = \Delta \cup T_{\text{STS}}$.

Fact

$M_F \models T_{\text{STS}}^*$.

There is more.

Theorem

*The theory T_{STS}^**

- *axiomatises the existentially closed Steiner triple systems*
- *it is model complete*
- *is the model companion of T_{STS}*
- *is complete*
- *has quantifier elimination.*

M_F is a prime model of T_{STS}^ .*

Let $T_{\text{STS}}^* = \Delta \cup T_{\text{STS}}$.

Fact

$M_F \models T_{\text{STS}}^*$.

There is more.

Theorem

*The theory T_{STS}^**

- *axiomatises the existentially closed Steiner triple systems*
- *it is model complete*
- *is the model companion of T_{STS}*
- *is complete*
- *has quantifier elimination.*

M_F is a prime model of T_{STS}^ .*

- T_{STS} is universal, so every model extends to an e.c. model
- T_{STS}^* axiomatises the e.c. models of T_{STS}

Therefore T_{STS}^* is the model companion of T_{STS} .

In particular, T_{STS}^* is model complete.

T_{STS}^* has the joint embedding property (because T_{STS} has), and it is model complete. Therefore T_{STS}^* is complete.

T_{STS}^* has the amalgamation property (because T_{STS} has), and it is model complete. Therefore T_{STS}^* has quantifier elimination.

- T_{STS} is universal, so every model extends to an e.c. model
- T_{STS}^* axiomatises the e.c. models of T_{STS}

Therefore T_{STS}^* is the model companion of T_{STS} .

In particular, T_{STS}^* is model complete.

T_{STS}^* has the joint embedding property (because T_{STS} has), and it is model complete. Therefore T_{STS}^* is complete.

T_{STS}^* has the amalgamation property (because T_{STS} has), and it is model complete. Therefore T_{STS}^* has quantifier elimination.

- T_{STS} is universal, so every model extends to an e.c. model
- T_{STS}^* axiomatises the e.c. models of T_{STS}

Therefore T_{STS}^* is the model companion of T_{STS} .

In particular, T_{STS}^* is model complete.

T_{STS}^* has the joint embedding property (because T_{STS} has), and it is model complete. Therefore T_{STS}^* is complete.

T_{STS}^* has the amalgamation property (because T_{STS} has), and it is model complete. Therefore T_{STS}^* has quantifier elimination.

- T_{STS} is universal, so every model extends to an e.c. model
- T_{STS}^* axiomatises the e.c. models of T_{STS}

Therefore T_{STS}^* is the model companion of T_{STS} .

In particular, T_{STS}^* is model complete.

T_{STS}^* has the joint embedding property (because T_{STS} has), and it is model complete. Therefore T_{STS}^* is complete.

T_{STS}^* has the amalgamation property (because T_{STS} has), and it is model complete. Therefore T_{STS}^* has quantifier elimination.

- T_{STS} is universal, so every model extends to an e.c. model
- T_{STS}^* axiomatises the e.c. models of T_{STS}

Therefore T_{STS}^* is the model companion of T_{STS} .

In particular, T_{STS}^* is model complete.

T_{STS}^* has the joint embedding property (because T_{STS} has), and it is model complete. Therefore T_{STS}^* is complete.

T_{STS}^* has the amalgamation property (because T_{STS} has), and it is model complete. Therefore T_{STS}^* has quantifier elimination.

The Fraïssé limit M_F is not saturated: it does not realise the type of a finitely generated infinite STS.

Question

Does T_{STS}^ have a countable saturated model?*

The Fraïssé limit M_F is not saturated: it does not realise the type of a finitely generated infinite STS.

Question

Does T_{STS}^ have a countable saturated model?*

A free construction

Given three points $\{a, b, c\}$ not forming a block, define a chain $\{A_i : i < \omega\}$ of partial STSs inductively:

- $A_0 = \{a, b, c\}$ and no products are defined
- $A_{i+1} = A_i \cup \{u \cdot v : u, v \in A_i \text{ and } u \cdot v \notin A_i\}$.

Then $A = \bigcup_{i < \omega} A_i$ is a countable STS.

A is the free STS on three generators (in the sense of universal algebra).

This construction can be generalised to include the disjoint union of partial STSs in the base step.

A free construction

Given three points $\{a, b, c\}$ not forming a block, define a chain $\{A_i : i < \omega\}$ of partial STSs inductively:

- $A_0 = \{a, b, c\}$ and no products are defined
- $A_{i+1} = A_i \cup \{u \cdot v : u, v \in A_i \text{ and } u \cdot v \notin A_i\}$.

Then $A = \bigcup_{i < \omega} A_i$ is a countable STS.

A is the free STS on three generators (in the sense of universal algebra).

This construction can be generalised to include the disjoint union of partial STSs in the base step.

A free construction

Given three points $\{a, b, c\}$ not forming a block, define a chain $\{A_i : i < \omega\}$ of partial STSs inductively:

- $A_0 = \{a, b, c\}$ and no products are defined
- $A_{i+1} = A_i \cup \{u \cdot v : u, v \in A_i \text{ and } u \cdot v \notin A_i\}$.

Then $A = \bigcup_{i < \omega} A_i$ is a countable STS.

A is the free STS on three generators (in the sense of universal algebra).

This construction can be generalised to include the disjoint union of partial STSs in the base step.

A free construction

Given three points $\{a, b, c\}$ not forming a block, define a chain $\{A_i : i < \omega\}$ of partial STSs inductively:

- $A_0 = \{a, b, c\}$ and no products are defined
- $A_{i+1} = A_i \cup \{u \cdot v : u, v \in A_i \text{ and } u \cdot v \notin A_i\}$.

Then $A = \bigcup_{i < \omega} A_i$ is a countable STS.

A is the free STS on three generators (in the sense of universal algebra).

This construction can be generalised to include the disjoint union of partial STSs in the base step.

A free construction

Given three points $\{a, b, c\}$ not forming a block, define a chain $\{A_i : i < \omega\}$ of partial STSs inductively:

- $A_0 = \{a, b, c\}$ and no products are defined
- $A_{i+1} = A_i \cup \{u \cdot v : u, v \in A_i \text{ and } u \cdot v \notin A_i\}$.

Then $A = \bigcup_{i < \omega} A_i$ is a countable STS.

A is the free STS on three generators (in the sense of universal algebra).

This construction can be generalised to include the disjoint union of partial STSs in the base step.

A free construction

Given three points $\{a, b, c\}$ not forming a block, define a chain $\{A_i : i < \omega\}$ of partial STSs inductively:

- $A_0 = \{a, b, c\}$ and no products are defined
- $A_{i+1} = A_i \cup \{u \cdot v : u, v \in A_i \text{ and } u \cdot v \notin A_i\}$.

Then $A = \bigcup_{i < \omega} A_i$ is a countable STS.

A is the free STS on three generators (in the sense of universal algebra).

This construction can be generalised to include the disjoint union of partial STSs in the base step.

Idea: we use the free construction to find 2^ω complete 3-types over \emptyset . This shows that T_{STS}^* does not have a countable saturated model.

Lemma

Let $\{A_i : i < \omega\}$ be a family of finite STSs such that $|A_i| \geq 3$ for at least one $i \in \omega$. Then there is a countably infinite M such that

- *M is generated by three elements*
- *every A_i embeds in M*
- *if a finite STS embeds in M , then it embeds in some A_i .*

Lemma (Doyen, 1969)

For all $n \equiv 1, 3 \pmod{6}$ there is an STS of cardinality n that does not embed any STS of cardinality m for $3 < m < n$.

Idea: we use the free construction to find 2^ω complete 3-types over \emptyset . This shows that T_{STS}^* does not have a countable saturated model.

Lemma

Let $\{A_i : i < \omega\}$ be a family of finite STSs such that $|A_i| \geq 3$ for at least one $i \in \omega$. Then there is a countably infinite M such that

- *M is generated by three elements*
- *every A_i embeds in M*
- *if a finite STS embeds in M , then it embeds in some A_i .*

Lemma (Doyen, 1969)

For all $n \equiv 1, 3 \pmod{6}$ there is an STS of cardinality n that does not embed any STS of cardinality m for $3 < m < n$.

Idea: we use the free construction to find 2^ω complete 3-types over \emptyset . This shows that T_{STS}^* does not have a countable saturated model.

Lemma

Let $\{A_i : i < \omega\}$ be a family of finite STSs such that $|A_i| \geq 3$ for at least one $i \in \omega$. Then there is a countably infinite M such that

- *M is generated by three elements*
- *every A_i embeds in M*
- *if a finite STS embeds in M , then it embeds in some A_i .*

Lemma (Doyen, 1969)

For all $n \equiv 1, 3 \pmod{6}$ there is an STS of cardinality n that does not embed any STS of cardinality m for $3 < m < n$.

Idea: we use the free construction to find 2^ω complete 3-types over \emptyset . This shows that T_{STS}^* does not have a countable saturated model.

Lemma

Let $\{A_i : i < \omega\}$ be a family of finite STSs such that $|A_i| \geq 3$ for at least one $i \in \omega$. Then there is a countably infinite M such that

- *M is generated by three elements*
- *every A_i embeds in M*
- *if a finite STS embeds in M , then it embeds in some A_i .*

Lemma (Doyen, 1969)

For all $n \equiv 1, 3 \pmod{6}$ there is an STS of cardinality n that does not embed any STS of cardinality m for $3 < m < n$.

Idea: we use the free construction to find 2^ω complete 3-types over \emptyset . This shows that T_{STS}^* does not have a countable saturated model.

Lemma

Let $\{A_i : i < \omega\}$ be a family of finite STSs such that $|A_i| \geq 3$ for at least one $i \in \omega$. Then there is a countably infinite M such that

- *M is generated by three elements*
- *every A_i embeds in M*
- *if a finite STS embeds in M , then it embeds in some A_i .*

Lemma (Doyen, 1969)

For all $n \equiv 1, 3 \pmod{6}$ there is an STS of cardinality n that does not embed any STS of cardinality m for $3 < m < n$.

Idea: we use the free construction to find 2^ω complete 3-types over \emptyset . This shows that T_{STS}^* does not have a countable saturated model.

Lemma

Let $\{A_i : i < \omega\}$ be a family of finite STSs such that $|A_i| \geq 3$ for at least one $i \in \omega$. Then there is a countably infinite M such that

- *M is generated by three elements*
- *every A_i embeds in M*
- *if a finite STS embeds in M , then it embeds in some A_i .*

Lemma (Doyen, 1969)

For all $n \equiv 1, 3 \pmod{6}$ there is an STS of cardinality n that does not embed any STS of cardinality m for $3 < m < n$.

Theorem

T_{STS}^* does not have a countable saturated model.

Proof.

(Sketch) For $n \equiv 1, 3 \pmod{6}$ let A_n be the STS of cardinality n given by Doyen's result.

Let $I = \{n \in \mathbb{N} : n \equiv 1, 3 \pmod{6}\}$.

For every infinite $X \subseteq I$ let M_X be the countable STS obtained by $\{A_n : n \in I\}$ as in the previous lemma. In particular, M_X is generated by three elements.

If $X \neq Y$, M_X and M_Y are not isomorphic. Then the types of the generators of M_X and M_Y are different.

Hence there are 2^ω of these types.



Theorem

T_{STS}^* does not have a countable saturated model.

Proof.

(Sketch) For $n \equiv 1, 3 \pmod{6}$ let A_n be the STS of cardinality n given by Doyen's result.

Let $I = \{n \in \mathbb{N} : n \equiv 1, 3 \pmod{6}\}$.

For every infinite $X \subseteq I$ let M_X be the countable STS obtained by $\{A_n : n \in I\}$ as in the previous lemma. In particular, M_X is generated by three elements.

If $X \neq Y$, M_X and M_Y are not isomorphic. Then the types of the generators of M_X and M_Y are different.

Hence there are 2^ω of these types.



Theorem

T_{STS}^* does not have a countable saturated model.

Proof.

(Sketch) For $n \equiv 1, 3 \pmod{6}$ let A_n be the STS of cardinality n given by Doyen's result.

Let $I = \{n \in \mathbb{N} : n \equiv 1, 3 \pmod{6}\}$.

For every infinite $X \subseteq I$ let M_X be the countable STS obtained by $\{A_n : n \in I\}$ as in the previous lemma. In particular, M_X is generated by three elements.

If $X \neq Y$, M_X and M_Y are not isomorphic. Then the types of the generators of M_X and M_Y are different.

Hence there are 2^ω of these types.



Theorem

T_{STS}^* does not have a countable saturated model.

Proof.

(Sketch) For $n \equiv 1, 3 \pmod{6}$ let A_n be the STS of cardinality n given by Doyen's result.

Let $I = \{n \in \mathbb{N} : n \equiv 1, 3 \pmod{6}\}$.

For every infinite $X \subseteq I$ let M_X be the countable STS obtained by $\{A_n : n \in I\}$ as in the previous lemma. In particular, M_X is generated by three elements.

If $X \neq Y$, M_X and M_Y are not isomorphic. Then the types of the generators of M_X and M_Y are different.

Hence there are 2^ω of these types.



Theorem

T_{STS}^* does not have a countable saturated model.

Proof.

(Sketch) For $n \equiv 1, 3 \pmod{6}$ let A_n be the STS of cardinality n given by Doyen's result.

Let $I = \{n \in \mathbb{N} : n \equiv 1, 3 \pmod{6}\}$.

For every infinite $X \subseteq I$ let M_X be the countable STS obtained by $\{A_n : n \in I\}$ as in the previous lemma. In particular, M_X is generated by three elements.

If $X \neq Y$, M_X and M_Y are not isomorphic. Then the types of the generators of M_X and M_Y are different.

Hence there are 2^ω of these types.



Theorem

T_{STS}^* does not have a countable saturated model.

Proof.

(Sketch) For $n \equiv 1, 3 \pmod{6}$ let A_n be the STS of cardinality n given by Doyen's result.

Let $I = \{n \in \mathbb{N} : n \equiv 1, 3 \pmod{6}\}$.

For every infinite $X \subseteq I$ let M_X be the countable STS obtained by $\{A_n : n \in I\}$ as in the previous lemma. In particular, M_X is generated by three elements.

If $X \neq Y$, M_X and M_Y are not isomorphic. Then the types of the generators of M_X and M_Y are different.

Hence there are 2^ω of these types.



Theorem

T_{STS}^* does not have a countable saturated model.

Proof.

(Sketch) For $n \equiv 1, 3 \pmod{6}$ let A_n be the STS of cardinality n given by Doyen's result.

Let $I = \{n \in \mathbb{N} : n \equiv 1, 3 \pmod{6}\}$.

For every infinite $X \subseteq I$ let M_X be the countable STS obtained by $\{A_n : n \in I\}$ as in the previous lemma. In particular, M_X is generated by three elements.

If $X \neq Y$, M_X and M_Y are not isomorphic. Then the types of the generators of M_X and M_Y are different.

Hence there are 2^ω of these types.



- T_{STS}^* is TP_2 and $NSOP_1$
- T_{STS}^* has elimination of hyperimaginaries and weak elimination of imaginaries.

- T_{STS}^* is TP_2 and $NSOP_1$
- T_{STS}^* has elimination of hyperimaginaries and weak elimination of imaginaries.